



Национальный исследовательский университет  
«Высшая школа экономики»

Международный центр анализа  
и выбора решений (DeCAn Lab)

Москва  
2022

# **Устойчивые турнирные решения как инструменты для принятия оптимальных решений: проблема обобщения двухпартийного множества на случай неполных предпочтений**

Андрей Субочев  
Ангелина Юдина



## Alternatives, comparisons, choices

$A$  – the *general set* of alternatives.

$X$  – the *menu*:  $X \subseteq A \wedge X \neq \emptyset \wedge |X| < \infty$ .

$R$  – results of binary comparisons,  $R \subseteq A \times A$ .

$R$  is presumed to be complete:  $\forall x \in A, \forall y \in A, (x, y) \in R \vee (y, x) \in R$ .

$R|_X = R \cap X \times X$  – restriction of  $R$  onto  $X$ .

$(X, R|_X)$  – *abstract game or weak tournament*.

$P$  – asymmetric part of  $R$ ,  $P \subseteq R$ :  $(x, y) \in P \Leftrightarrow ((x, y) \in R \wedge (y, x) \notin R)$ .

If  $P|_X$  is complete,  $\forall x \in A, \forall y \in A \wedge y \neq x, (x, y) \in P \vee (y, x) \in P$ , then

$(X, R|_X)$  – *(proper) tournament*.



## Tournament solutions

A tournament solution  $S$  is a correspondence

$$S(X, R): 2^A \setminus \emptyset \times 2^{A \times A} \rightarrow 2^A$$

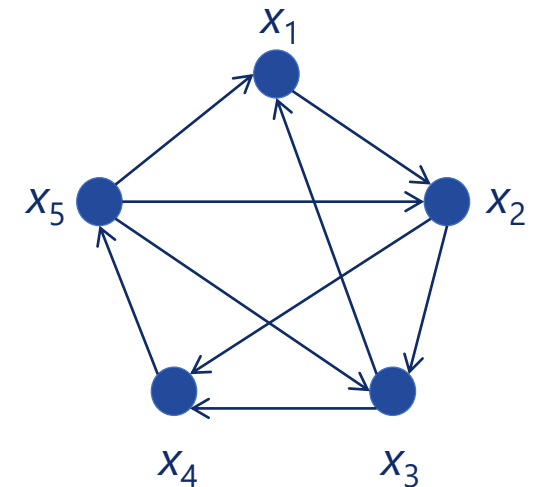
that has the following properties:

0. *Locality*:  $S(X, R) = S(R|_X) \subseteq X$
1. *Nonemptiness*:  $\forall X, \forall R, S(R|_X) \neq \emptyset$ ;
2. *Neutrality*: permutation of alternatives' names and choice commute;
3. *Condorcet consistency*:

$$\text{MAX}(R|_X) \subseteq S(R|_X) \wedge \text{MAX}(R|_X) = \{w\} \Rightarrow S(R|_X) = \{w\}.$$

	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$
$x_1$	0	1	0	0	0
$x_2$	0	0	1	1	0
$x_3$	1	0	0	1	0
$x_4$	0	0	0	0	1
$x_5$	1	1	1	0	0

Tournament matrix  $T$



Tournament digraph



## Lotteries

Comparison function:  $g(x_1, x_2)=1 \Leftrightarrow x_1 P x_2$ ,  $g(x_1, x_2)=-1 \Leftrightarrow x_2 P x_1$ , otherwise  $g(x_1, x_2)=0$ .

Since matrix  $G = \|g(x_i, x_j)\|$  is skew-symmetric,

formula  $p_1 G p_2$  defines a binary relation on the set of lotteries:  $p_1 G p_2 \geq 0 \Leftrightarrow p_1 \succsim p_2$ .

If  $p_0 G p \geq 0$  for all  $p$  then  $p_0$  is a *maximal lottery*.

The set  $\{x\}$  is the support of a maximal lottery on  $X \Leftrightarrow x$  is a maximal element of  $R|_X$ .

	$x_1$	$x_2$	$x_3$
$x_1$	0	1	0
$x_2$	0	0	1
$x_3$	1	0	0

Tournament matrix T

	$x_1$	$x_2$	$x_3$
$x_1$	0	1	-1
$x_2$	-1	0	1
$x_3$	1	-1	0

Matrix G



## Bipartisan set ( $BP$ ) and Essential set ( $E$ )

1. The set of maximal lotteries is always nonempty.
2. If a tournament  $(X, R|_X)$  is proper then there is just one maximal lottery.

***Bipartisan set  $BP$***  (Laffond, Laslier, Le Breton, 1993)

of a (proper) tournament  $(X, R|_X)$  is the support of the (unique) maximal lottery.

***Essential set  $E$***  (Dutta, Laslier, 1999)

of a (weak) tournament  $(X, R|_X)$  is the union of supports of all maximal lotteries.



## Example

The Condorcet cycle.

$$X = \{x_1, x_2, x_3\}, R|_X = \{(x_1, x_2), (x_2, x_3), (x_3, x_1)\}.$$

Maximal lottery  $\mathbf{p}_{\max} = (1/3, 1/3, 1/3)$ .

Bipartisan set  $BP = X$ .

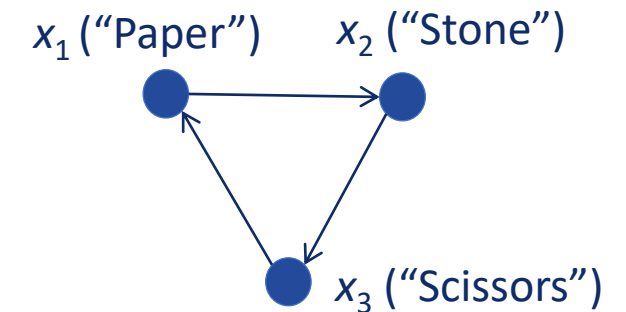
Note that  $\mathbf{p}_{\max}$  is an eigenvector of  $\mathbf{G}$  with the eigenvalue 0,  
therefore  $\mathbf{p}\mathbf{G}\mathbf{p}_{\max} = 0$  for all  $\mathbf{p}$ .

	$x_1$	$x_2$	$x_3$
$x_1$	0	1	0
$x_2$	0	0	1
$x_3$	1	0	0

Tournament matrix  $\mathbf{T}$

	$x_1$	$x_2$	$x_3$
$x_1$	0	1	-1
$x_2$	-1	0	1
$x_3$	1	-1	0

Matrix  $\mathbf{G}$



Tournament game –  
"Paper, Scissors, Stone"



## Properties

- **Monotonicity**

$$\forall R_1, R_2 \subseteq A^2, \forall X \subseteq A, \forall x \in S(R_1|_X),$$

$$(R_1|_{X \setminus \{x\}} = R_2|_{X \setminus \{x\}} \wedge \forall y \in X \setminus \{x\}, (xP_1y \Rightarrow xP_2y) \wedge (xR_1y \Rightarrow xR_2y)) \Rightarrow x \in S(R_2|_X).$$

- **Stability**

For all  $R \subseteq A^2$  and for all  $X, Y \subseteq A$  such that  $X \cap Y \neq \emptyset$  the following holds:

$$S(X, R) = S(Y, R) = Z \Leftrightarrow S(X \cup Y, R) = Z \wedge Z \subseteq X \cap Y.$$

- **Computational simplicity**

There is a polynomial algorithm for computing  $S$ .



## Properties related to stability

**Stability:**  $S(X, R)=S(Y, R)=Z \Leftrightarrow S(X \cup Y, R)=Z \wedge Z \subseteq X \cap Y$ .

- **$\alpha$ -property** (*Outcast property or generalized Nash independence of irrelevant alternatives*):

$$S(X, R)=S(Y, R)=Z \Leftarrow S(X \cup Y, R)=Z \wedge Z \subseteq X \cap Y.$$

- **$\gamma$ -property:**

$$S(X, R)=S(Y, R)=Z \Rightarrow S(X \cup Y, R)=Z \wedge Z \subseteq X \cap Y.$$

- **Idempotence:**  $\forall X, S(S(X))=S(X)$ .

- **The Aizerman-Aleskerov condition:**  $\forall X, \forall Y, S(X) \subseteq Y \subseteq X \Rightarrow S(Y) \subseteq S(X)$ .

- **Independence of irrelevant results (independence of losers):**

$$\forall R_1, R_2 \subseteq A^2, \forall X \subseteq A, (\forall x \in S(R_1|_X), \forall y \in X, ((xR_1y \Leftrightarrow xR_2y) \wedge (yR_1x \Leftrightarrow yR_2x))) \Rightarrow S(R_1|_X)=S(R_2|_X).$$

$\alpha$ -property  $\Leftrightarrow$  Idempotence  $\wedge$  the Aizerman-Aleskerov condition

$\alpha$ -property  $\Rightarrow$  Independence of irrelevant results





## The conservative extension (Brandt et al., 2014, 2018)

A tournament  $(X, T)$  is called **orientation** of a weak tournament  $(X, P)$  if  $(X, T)$  is proper and  $P \subseteq T$ .

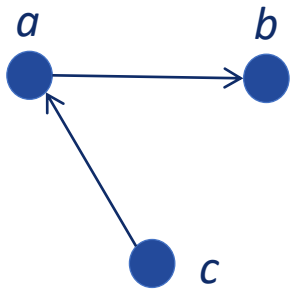
For a tournament solution  $S(X, P)$ , its **conservative extension** (denoted  $[S]$ ) to weak tournaments is the choice correspondence  $[S](X, P)$  defined thus:

an alternative  $x$  from  $X$  belongs to  $[S](X, P)$  if and only if there is an **orientation**  $(X, T)$  of  $(X, P)$ , such that  $x$  belong to  $S(X, T)$ .

**Theorem:** The conservative extension preserves properties of the original solution.

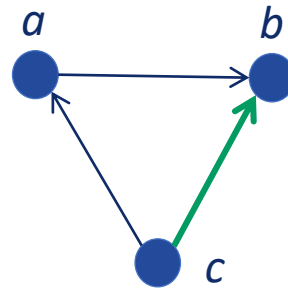


## Example



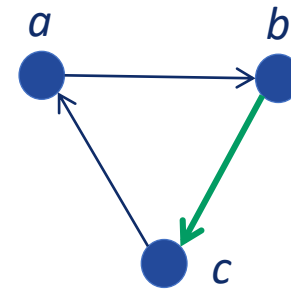
Weak tournament digraph

$[BP]=?$

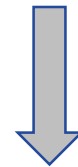


Orientations of the weak tournament

$BP=\{c\}$



$BP=\{a, b, c\}$



$[BP]=\{a, b, c\}$



## Axiomatic analysis

	<b><i>BP</i></b>	<b><i>[BP]</i></b>	<b><i>E</i></b>
Monotonicity	Yes	Yes	Yes
$\alpha$ -property (outcast)	Yes	Yes	Yes
Idempotence	Yes	Yes	Yes
Aizerman-Aleskerov property	Yes	Yes	Yes
Independence of irrelevant results	Yes	Yes	Yes
$\gamma$ -property	Yes	Yes	Yes
Stability	Yes	Yes	Yes
Computational simplicity	Yes	Yes	Yes



## Relations of $E$ and $[BP]$ to other solutions

In proper tournaments,  $E=[BP]=BP \subseteq UC \subseteq ES$ , also  $BP \subseteq D \subseteq ES$ .

In weak tournaments,

	$E$	$[BP]$	$UC_{IM}$	$UC_M$	$UC_{IF}$	$UC_F$	$UC_{McK}$	$UC_D$	$D$	$SP$	$[D]$	$WS$	$ES$	$RES$	$[ES]$	$UCp$	$STC$	$UT$	$WTC$
$E$	=	$\cap$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$	$\subset$	$\subset$	$\emptyset$	$\cap$	$\subset$	$\cap$	$\cap$	$\cap$	$\subset$	$\subset$	$\cap$	$\cap$	$\subset$
$[BP]$	$\cap$	=	$\cap$	$\cap$	$\cap$	$\cap$	$\cap$	$\subset$	$\cap$	$\cap$	$\subset$	$\cap$	$\cap$	$\cap$	$\subset$	$\cap$	$\cap$	$\cap$	$\subset$

If  $S_1$  and  $S_2$  are the tournament solutions denoting, correspondingly, a row and a column, then a symbol in the corresponding box means the following:

“ $\emptyset$ ” –  $S_1$  and  $S_2$  are independent, that is, their intersection can be empty in some finite tournament;

“ $\cap$ ” –  $S_1$  and  $S_2$  are not logically nested, but their intersection is always nonempty;

“ $\subset$ ” –  $S_1$  refines  $S_2$ ; “=” –  $S_1$  and  $S_2$  are identical.



$E$  and  $[BP]$  are not logically nested (Brandt et al., 2018)

We proved that  $E$  and  $[BP]$  are **not** independent:  $E \cap [BP] \neq \emptyset$  always.

**Lemma** (Tucker 1956): For any skew-symmetric matrix  $\mathbf{G}$  there exists a vector  $\mathbf{p}$  such that  $\mathbf{p} \geq \mathbf{0}$  and  $\mathbf{G}\mathbf{p} \geq \mathbf{0}$  and  $\mathbf{p}\mathbf{G}\mathbf{p} = 0$  and  $\mathbf{p} + \mathbf{G}\mathbf{p} > \mathbf{0}$ .

### Restatement of Tucker's lemma

For any antisymmetric comparison function  $g(x, y): X \times X \rightarrow \mathbb{R}$  there exists a lottery  $\mathbf{p}$  on  $X$  such that  $\forall y \in X, \sum_{x \in X} p(x)g(x, y) \geq 0$

and exactly one of the two numbers  $p(y)$  and  $\sum_{x \in X} p(x)g(x, y)$  is positive, while the other is 0.



Национальный исследовательский университет  
«Высшая школа экономики»

Международный центр анализа  
и выбора решений (DeCAn Lab)

Москва  
2022

# Спасибо за внимание!

101000, Россия,  
Москва, Мясницкая, 20  
тел: 495 621-7983  
факс: 495 628-7931  
[www.hse.ru](http://www.hse.ru)



## The covering relations and the uncovered sets

**The covering relations** (Fishburn, 1977; Miller, 1980; McKelvey, 1986; Duggan, 2007, 2013)

The covering relation  $C \subseteq X^2$ , is a strengthening of  $P|_X$ :

1. The Miller covering  $C_M: xC_M y \Leftrightarrow xPy \wedge P^{-1}(y) \subset P^{-1}(x)$ .
2. The weak Miller covering  $C_{WM}: xC_{WM} y \Leftrightarrow P^{-1}(y) \subset P^{-1}(x)$ .
3. The Fishburn covering  $C_F: xC_F y \Leftrightarrow xPy \wedge P(x) \subset P(y)$ .
4. The weak Fishburn covering  $C_{WF}: xC_{WF} y \Leftrightarrow P(x) \subset P(y)$ .
5. The McKelvey covering  $C_{MCK}: xC_{MCK} y \Leftrightarrow xPy \wedge P^{-1}(y) \subset P^{-1}(x) \wedge P(x) \subset P(y)$ .
6. The Duggan covering  $C_D: xC_D y \Leftrightarrow P^0(y) \cup P^{-1}(y) \subset P^{-1}(x)$ .

The set of all alternatives that are not (weakly) covered in  $X$  by any alternative is called **the (inner) uncovered set** of a feasible set  $X$ .

The Miller, Fishburn, McKelvey and Duggan uncovered sets and their inner versions will be denoted  $UC_M, UC_F, UC_{MCK}, UC_D, UC_{IM}$  and  $UC_{IF}$ , correspondingly.



## Minimal externally stable sets

A nonempty subset  $Y$  of  $X$  is called

*P-dominating* if  $\forall x \in X, \exists y \in Y: yPx$

*P-externally stable* if  $\forall x \in X \setminus Y, \exists y \in Y: yPx$

*R-externally stable* if  $\forall x \in X \setminus Y, \exists y \in Y: yRx$

*Self-protecting* if  $\forall x \in X, (\exists y \in Y: yPx) \vee (\forall y \in Y, yRx)$

*Weakly stable* if  $\forall x \in X \setminus Y, (\exists y \in Y: yPx) \vee (\forall y \in Y, yRx)$

**Tournament solutions:** the union of all minimal  
*P-dominating* sets  $D$  (Duggan 2013, Subochev 2016)  
*P-externally stable* sets  $ES$  (Wufl, Feld, Owen &  
Grofman 1989, Subochev 2008)  
*R-externally stable* sets  $RES$  (Aleskerov & Subochev  
2009, 2013)  
*Self-protecting* sets  $SP$  (Roth 1976, Subochev 2020)  
*Weakly stable* sets  $WS$  (Aleskerov & Kurbanov 1999)